

Note 4

# 04 Convex Optimization

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The goal in optimization is usually to find some parameters  $\theta$  that minimize a continuous **objective function**  $f(\theta)$ . Note  $f$  is considered **smooth** if its gradient is also continuous.

## 1 Convexity

Consider a **convex set**  $C$ , which means that for any points in  $C$  we select, the line between them will lie completely in  $C$ . How is this related to, say, a convex function? As it turns out, we can relate the two by considering the **epigraph** of a function  $f$ , the set of all points that lie above  $f$ .

### 1.1 Epigraph

First, let us formalize the definition of a **strict epigraph**. (The standard epigraph has a weak inequality):

$$\{(x, t) : x \in \mathbb{R}^n, t \in \mathbb{R}, t > f(x)\}$$

Second, consider the formal definition of a **strongly convex** function  $f$ . A function  $f$  is convex if between any two points on  $f$ , the line between them lies entirely above the function - formally, if it satisfies the following, for  $t \in [0, 1]$  and all  $x_1, x_2$ . Without loss of generality, consider  $x_1 < x_2$ .

$$f((1-t)x_1 + tx_2) < (1-t)f(x_1) + tf(x_2)$$

To define a **weakly convex** function, we simply convert the strict inequality into a weak one, to get  $f((1-t)x_1 + tx_2) \leq (1-t)f(x_1) + tf(x_2)$ .

**Explanation of Definition:**

Intuitively, we take the points  $(x_1, f(x_1)), (x_2, f(x_2))$  and draw the line between them. Consider taking your finger, and following the line from  $x_1$  to  $x_2$ .  $t$  represents how far along we are, on that line, where  $t = 0$  is at  $x_1$  and  $t = 1$  is at  $x_2$ . To compute the y-value at a certain point along that line, we simply weight the two y-values at the end points. If we are closer to  $x_1$ , weight  $f(x_1)$  more heavily. Thus, the y value on that line is

$$(1 - t)f(x_1) + tf(x_2)$$

The corresponding evaluation at the real value of  $f$  is straightforward. Instead of weighting the y values, we weight the x values, then evaluate the true y value at that point.

$$f((1 - t)x_1 + tx_2)$$

On an intuitive level, this is how the above definition captures the condition that any line lies above  $f$ , or in words, any line between two points on  $f$  will necessarily lie in the epigraph of  $f$ .

**Theorem:** The epigraph of a function  $f$  is a convex set if and only if  $f$  is weakly convex.

**Proof:**

The forward direction is simpler. Take a pair of points  $(x_i, f(x_i))$ . Since  $f$  is a convex set, we know for any point along the line between them,  $(x, y)$  must satisfy the condition  $y \geq f(x)$ . Plugging in for  $x = (1 - t)x_1 + tx_2$  and  $y = (1 - t)y_1 + ty_2 = (1 - t)f(x_1) + tf(x_2)$ , our claim immediately follows.

We then prove the reverse direction: if  $f$  is convex, the epigraph is a convex set. Take any two points  $(x_i, y_i)$  and the line between them.

$$x = (1 - t)x_1 + tx_2$$

$$y = (1 - t)y_1 + ty_2$$

Apply the definition of an epigraph, to plug in for  $y_i$ . Then, apply the definition of convexity.

$$y = (1 - t)y_1 + ty_2 \geq (1 - t)f(x_1) + tf(x_2) \geq f((1 - t)x_1 + tx_2) = f(x)$$

Thus,  $y \geq f(x)$  and  $(x, y)$  is in the epigraph.

## 1.2 Solutions

If we minimize over a convex function  $f$ , we could find a **local minimum**, a value  $x$  such that there exists a ball of some radius around  $x$  where all  $x'$  in this ball satisfy  $f(x') \geq f(x)$ . It may be a **global minimum**, where for *all*  $x'$ ,  $f(x') \geq f(x)$ . However, the global minimum is often difficult and sometimes impossible to find. A continuous convex function has:

1. No minimum, as the function may be weakly convex and go to  $-\infty$ .
2. Just one local minimum
3. Connected set of local minima that are all global minima.

Perceptrons fall in the last category, and neural networks often find local not global minima.

## 1.3 Cases

Consider the following cases when functions are convex.

1. Constant functions:  $f(x) = c$
2. Powers of x:  $f(x) = x^r$  on interval  $[0, \infty]$  when  $r \geq 1$
3. Affine transform:  $f(x)$  convex implies  $f(w^T x + \beta)$  is convex.
4. Log:  $-\log(x)$  is convex.
5. Reverse of Concave: If  $f(x)$  is concave,  $-f(x)$  is convex.
6. Combination: If  $f, g$  are convex, then  $\max(f(x), g(x))$  is convex.

## 2 Optimization Problems

### 2.1 Unconstrained Optimization

For smooth functions, we can use gradient descent, the nonlinear conjugate gradient, or Newton's method. For nonsmooth functions, we can use gradient descent. A **line search** finds a local minimum along one dimension.

### 2.2 Constrained Optimization

Here we minimize  $f(\theta)$  for some parameters  $\theta$  subject to  $g(\theta) = 0$  for some constraint  $g$ , where  $g$  is a smooth function. In other words, we have smooth equality constraints. To solve this, use Lagrange multipliers.

#### 2.2.1 Lagrange Multipliers

Lagrange multipliers allow us to solve optimization problems with constraints.

1. Set up a new function  $h(\theta) = f(\theta) + \lambda g(w)$ , where  $\lambda$  is the multiplier.
2. Take partials with respect to every parameter in  $\theta$  and  $\lambda$ . If there are  $d$  parameters, there are  $d + 1$  unknowns.
3. Set each partial to 0. This makes  $d + 1$  equations for  $d + 1$  unknowns.
4. Solve the system.
5. Ensure that you do not divide variables  $x_i$  out without considering the possibility that  $x_i = 0$ .

### 2.3 Linear Programming

With a linear objective function with inequality constraints, our goal is to maximize or minimize  $c \cdot w$  subject to  $Aw \leq b$ . Note the polytope or feasible region is convex. For some constraints - named the **active constraints** - the optimum achieves equality. In an SVM, vectors that likewise touch or cross the slab are called **support vectors**.

### 2.4 Quadratic Programming

Minimize  $f(w) = w^T Q w + c^T w$  subject to  $Aw \leq b$ , where  $Q$  is a symmetric, positive definite matrix.